

Figure 3.1 Network types

	<i>Range</i>	<i>Bandwidth (Mbps)</i>	<i>Latency (ms)</i>
LAN	1–2 kms	10–1000	1–10
WAN	worldwide	0.010–600	100–500
MAN	2–50 kms	1–150	10
Wireless LAN	0.15–1.5 km	2–11	5–20
Wireless WAN	worldwide	0.010–2	100–500
Internet	worldwide	0.010–2	100–500

Figure 3.2 Conceptual layering of protocol software

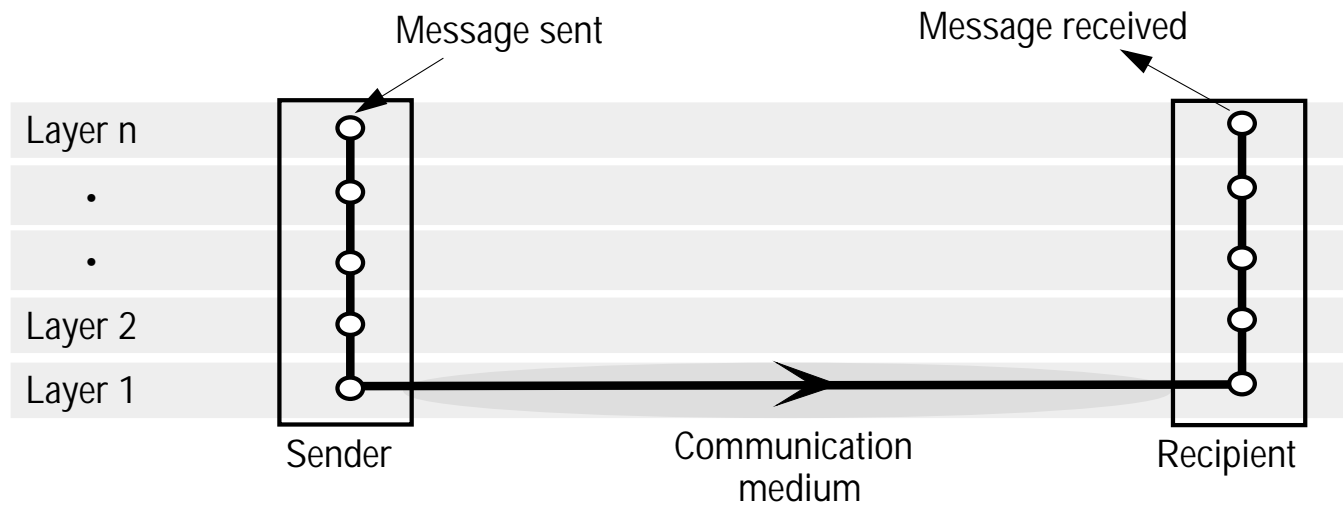


Figure 3.3 Encapsulation as it is applied in layered protocols

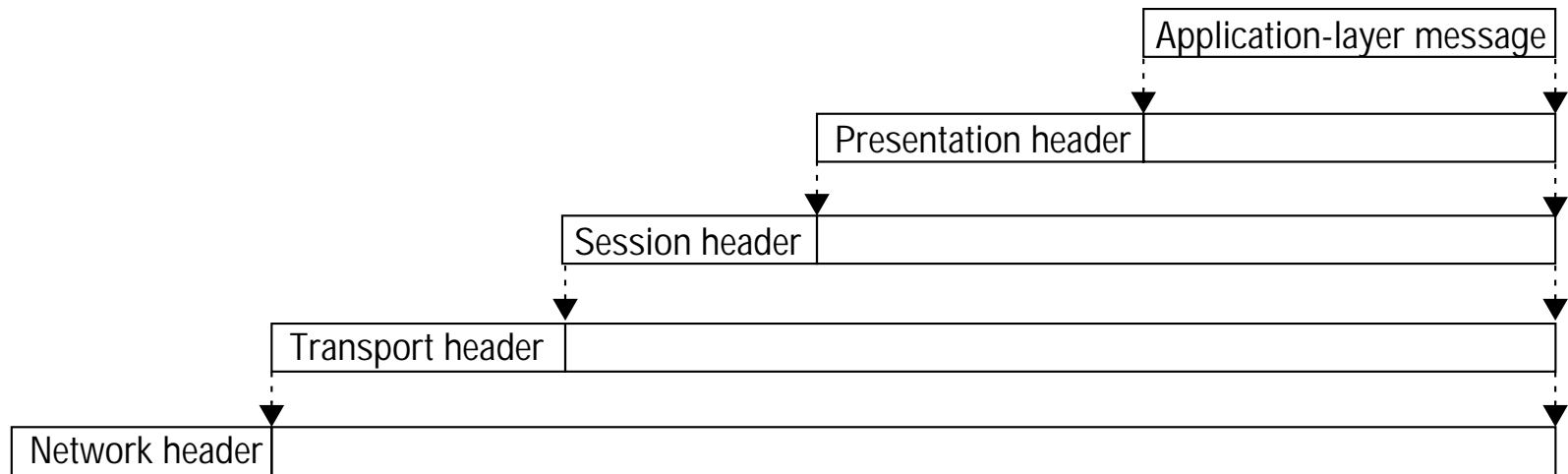


Figure 3.4 Protocol layers in the ISO *Open Systems Interconnection (OSI)* protocol model

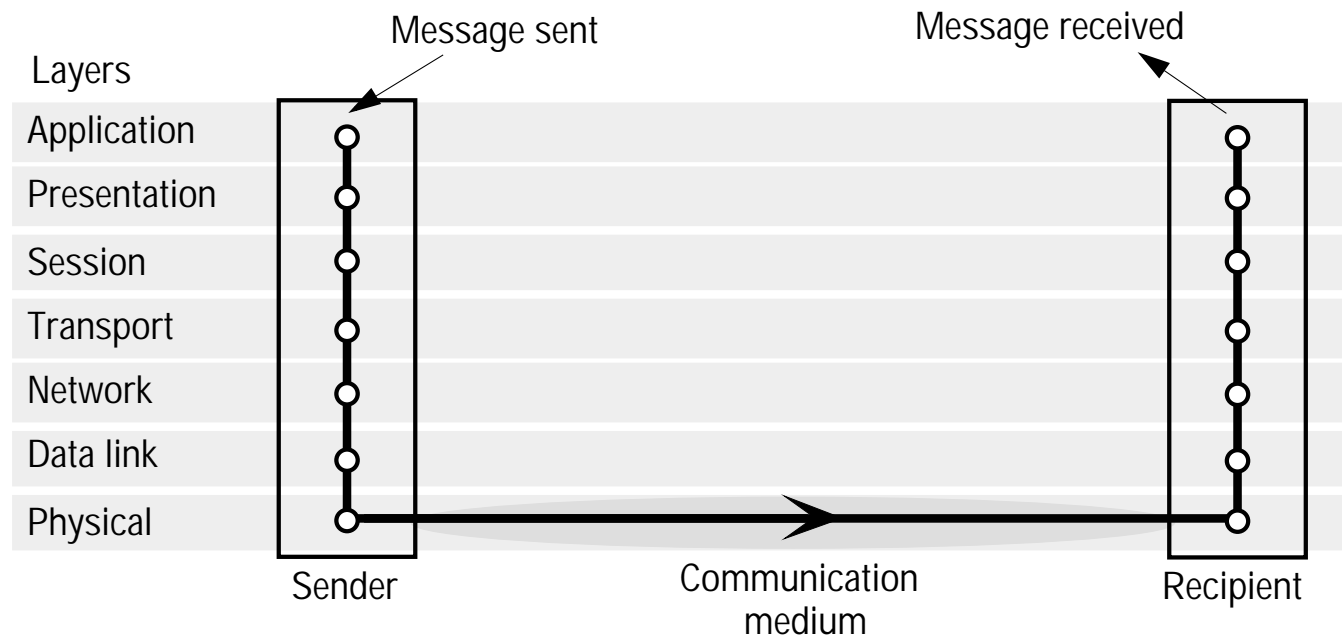


Figure 3.5 OSI protocol summary

<i>Layer</i>	<i>Description</i>	<i>Examples</i>
Application	Protocols that are designed to meet the communication requirements of specific applications, often defining the interface to a service.	HTTP, FTP, SMTP, CORBA IIOP
Presentation	Protocols at this level transmit data in a network representation that is independent of the representations used in individual computers, which may differ. Encryption is also performed in this layer, if required.	Secure Sockets (SSL), CORBA Data Rep.
Session	At this level reliability and adaptation are performed, such as detection of failures and automatic recovery.	
Transport	This is the lowest level at which messages (rather than packets) are handled. Messages are addressed to communication ports attached to processes. Protocols in this layer may be connection-oriented or connectionless.	TCP, UDP
Network	Transfers data packets between computers in a specific network. In a WAN or an internetwork this involves the generation of a route passing through routers. In a single LAN no routing is required.	IP, ATM virtual circuits
Data link	Responsible for transmission of packets between nodes that are directly connected by a physical link. In a WAN transmission is between pairs of routers or between routers and hosts. In a LAN it is between any pair of hosts.	Ethernet MAC, ATM cell transfer, PPP
Physical	The circuits and hardware that drive the network. It transmits sequences of binary data by analogue signalling, using amplitude or frequency modulation of electrical signals (on cable circuits), light signals (on fibre optic circuits) or other electromagnetic signals (on radio and microwave circuits).	Ethernet base-band signalling, ISDN

Figure 3.6 Internetwork layers

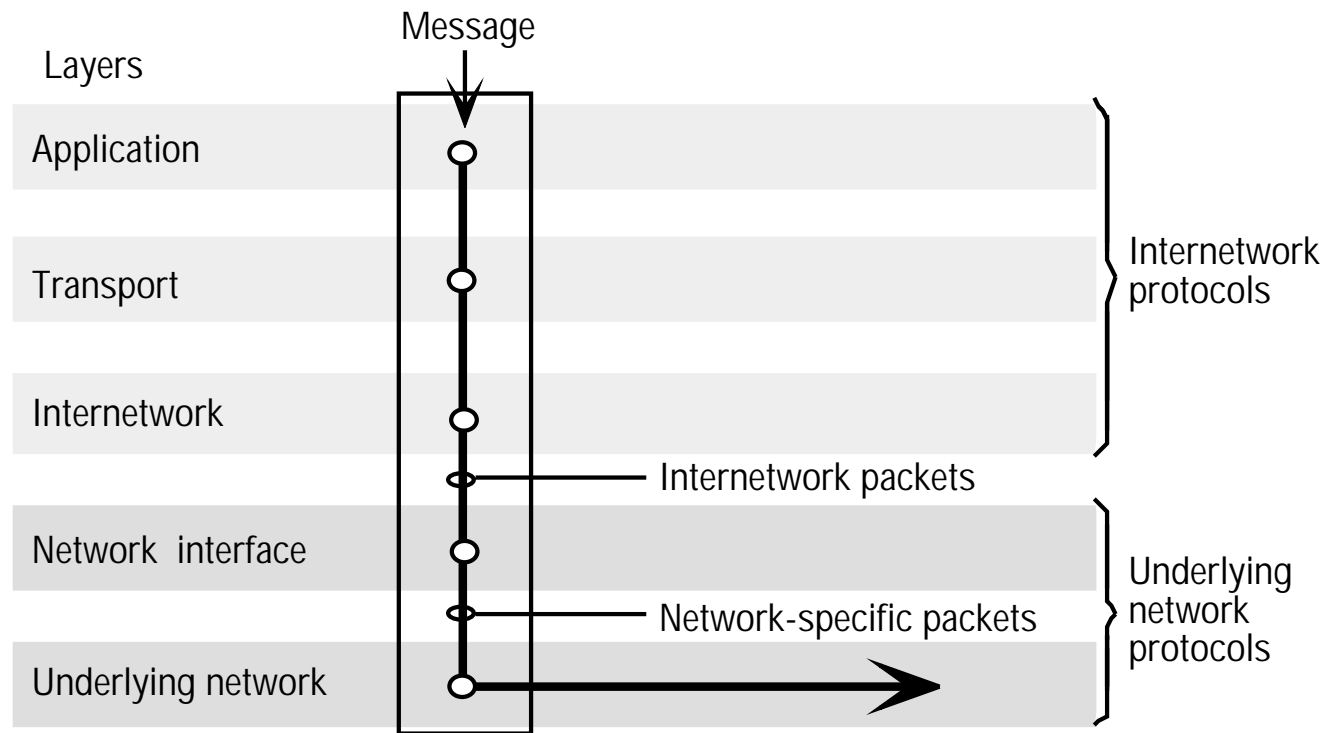


Figure 3.7 Routing in a wide area network

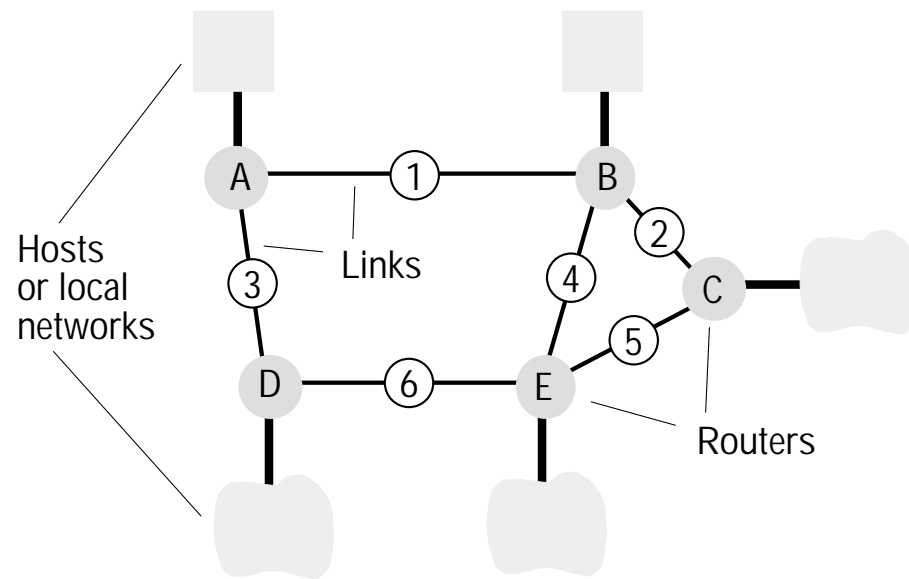


Figure 3.8 Routing tables for the network in Figure 3.7

<i>Routings from A</i>			<i>Routings from B</i>			<i>Routings from C</i>		
<i>To</i>	<i>Link</i>	<i>Cost</i>	<i>To</i>	<i>Link</i>	<i>Cost</i>	<i>To</i>	<i>Link</i>	<i>Cost</i>
A	local	0	A	1	1	A	2	2
B	1	1	B	local	0	B	2	1
C	1	2	C	2	1	C	local	0
D	3	1	D	1	2	D	5	2
E	1	2	E	4	1	E	5	1

<i>Routings from D</i>			<i>Routings from E</i>		
<i>To</i>	<i>Link</i>	<i>Cost</i>	<i>To</i>	<i>Link</i>	<i>Cost</i>
A	3	1	A	4	2
B	3	2	B	4	1
C	6	2	C	5	1
D	local	0	D	6	1
E	6	1	E	local	0

Figure 3.9 Pseudo-code for RIP routing algorithm

Send: Each t seconds or when Tl changes, send Tl on each non-faulty outgoing link.

Receive: Whenever a routing table Tr is received on link n :

```
for all rows  $Rr$  in  $Tr$  {
```

```
  if ( $Rr.link \neq n$ ) {
```

```
     $Rr.cost = Rr.cost + 1$ ;
```

```
     $Rr.link = n$ ;
```

```
    if ( $Rr.destination$  is not in  $Tl$ ) add  $Rr$  to  $Tl$ ; // add new destination to  $Tl$ 
```

```
  else for all rows  $Rl$  in  $Tl$  {
```

```
    if ( $Rr.destination = Rl.destination$  and
```

```
        ( $Rr.cost < Rl.cost$  or  $Rl.link = n$ ))  $Rl = Rr$ ;
```

```
        //  $Rr.cost < Rl.cost$  : remote node has better route
```

```
        //  $Rl.link = n$  : remote node is more authoritative
```

```
    }
```

```
  }
```

```
}
```

Figure 3.10 Simplified view of the QMW Computer Science network

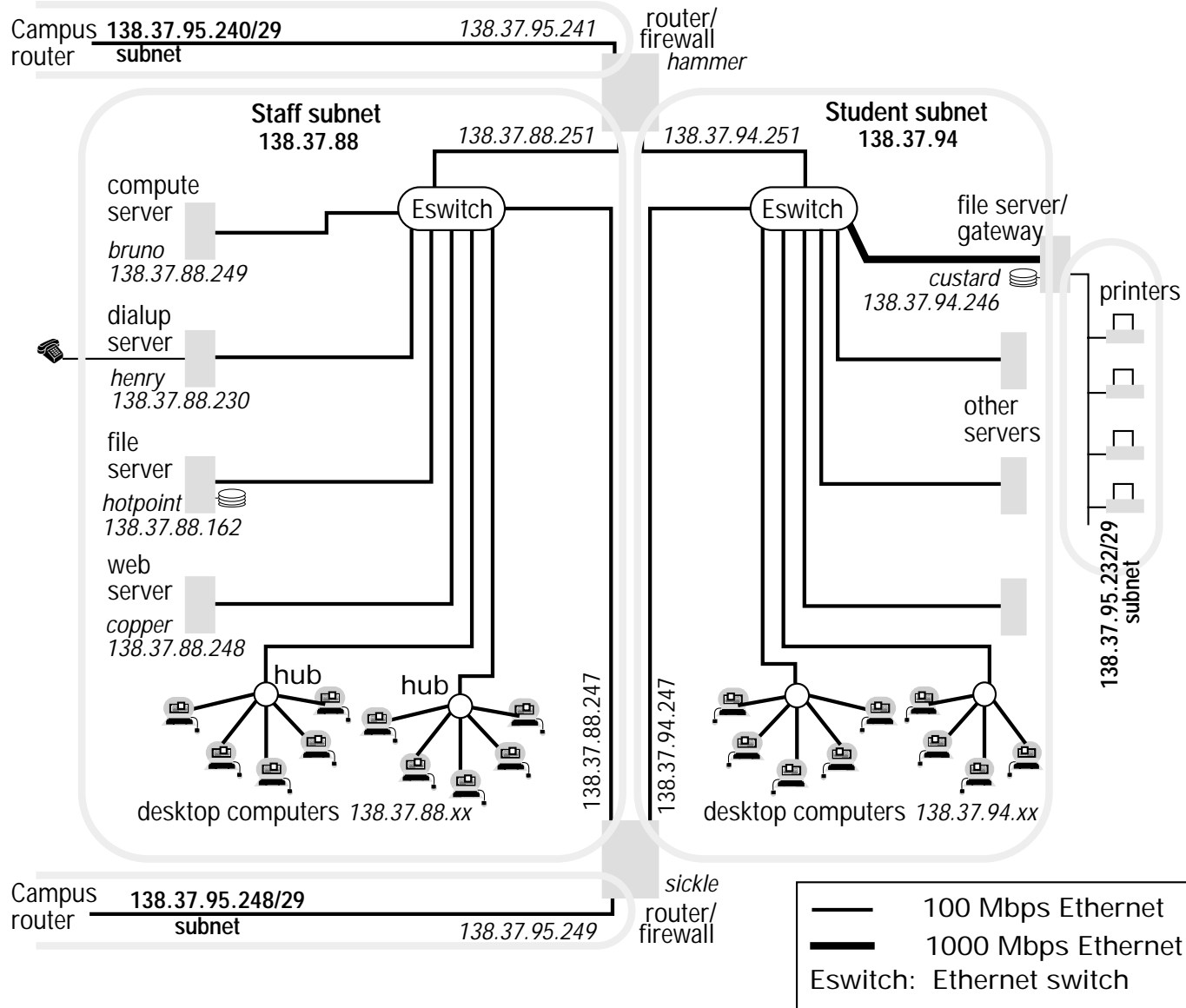


Figure 3.11 Tunnelling for IPv6 migration

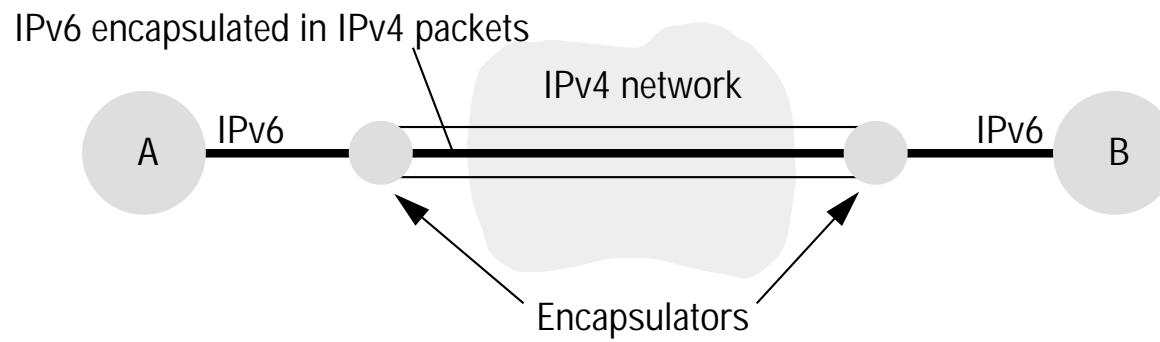


Figure 3.12 TCP/IP layers

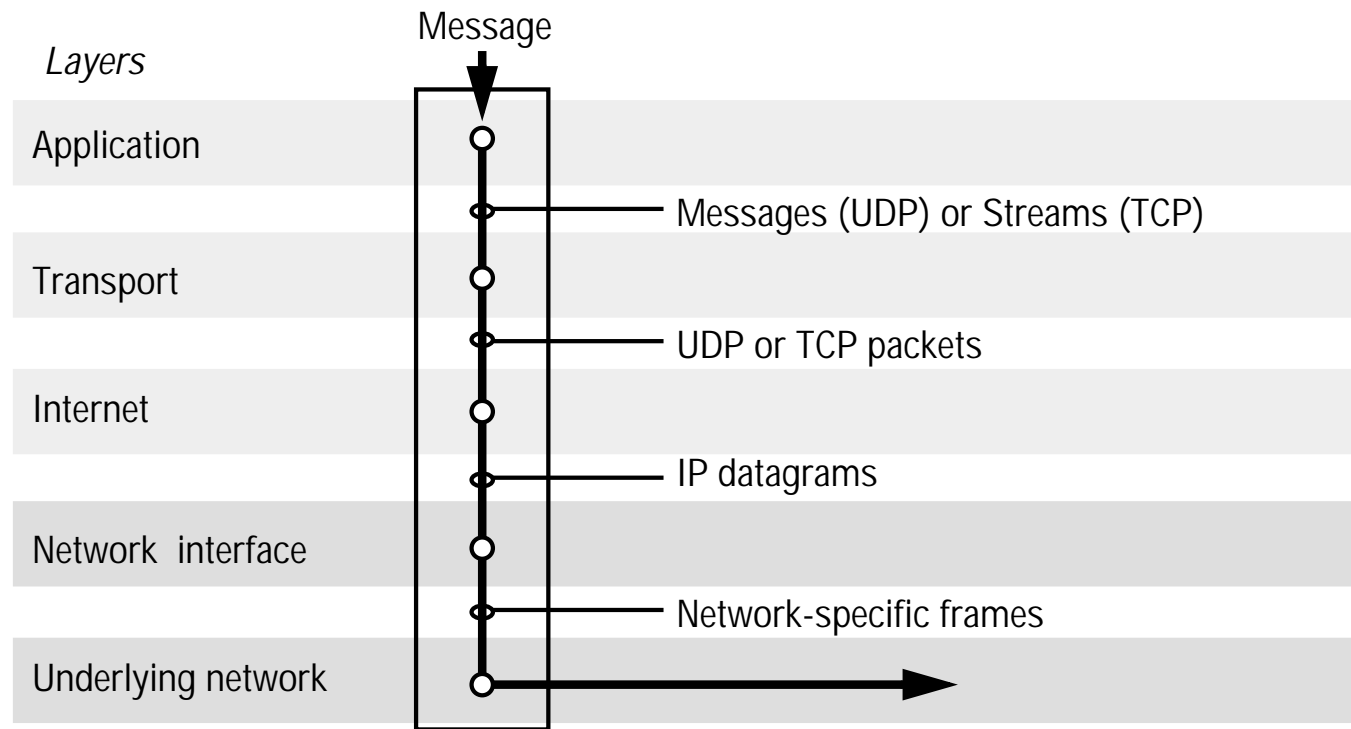


Figure 3.13 Encapsulation as it occurs when a message is transmitted via TCP over an Ethernet

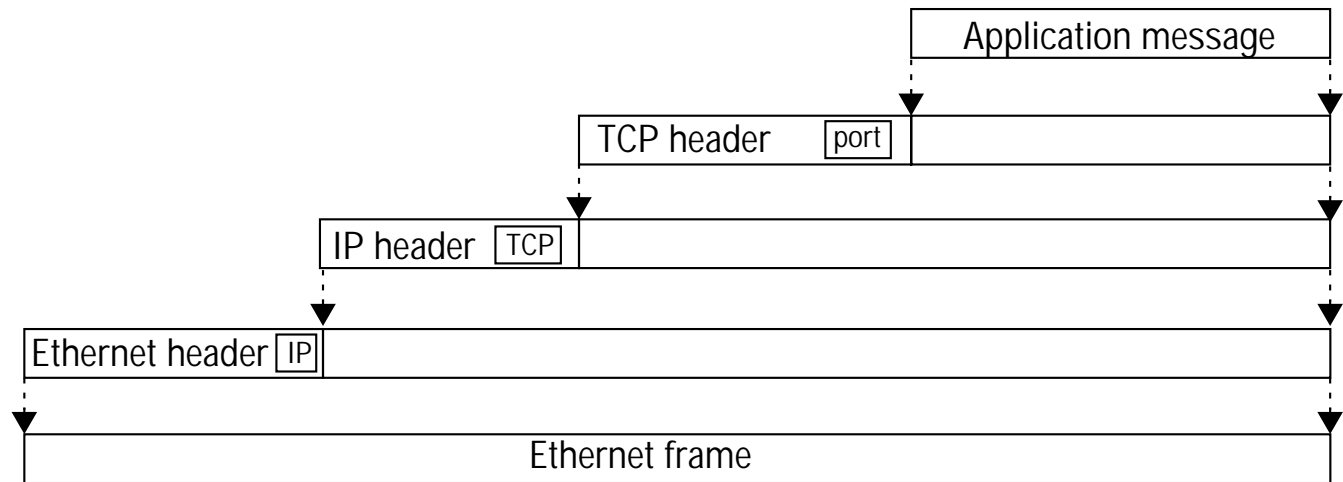


Figure 3.14 The programmer's conceptual view of a TCP/IP Internet

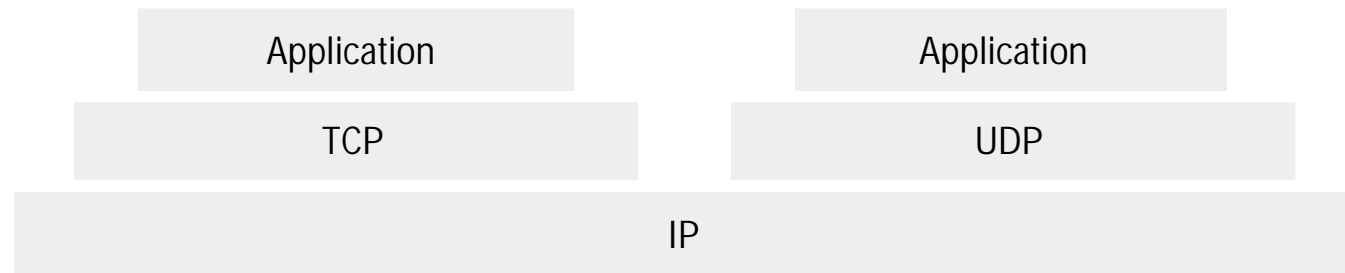


Figure 3.15 Internet address structure, showing field sizes in bits

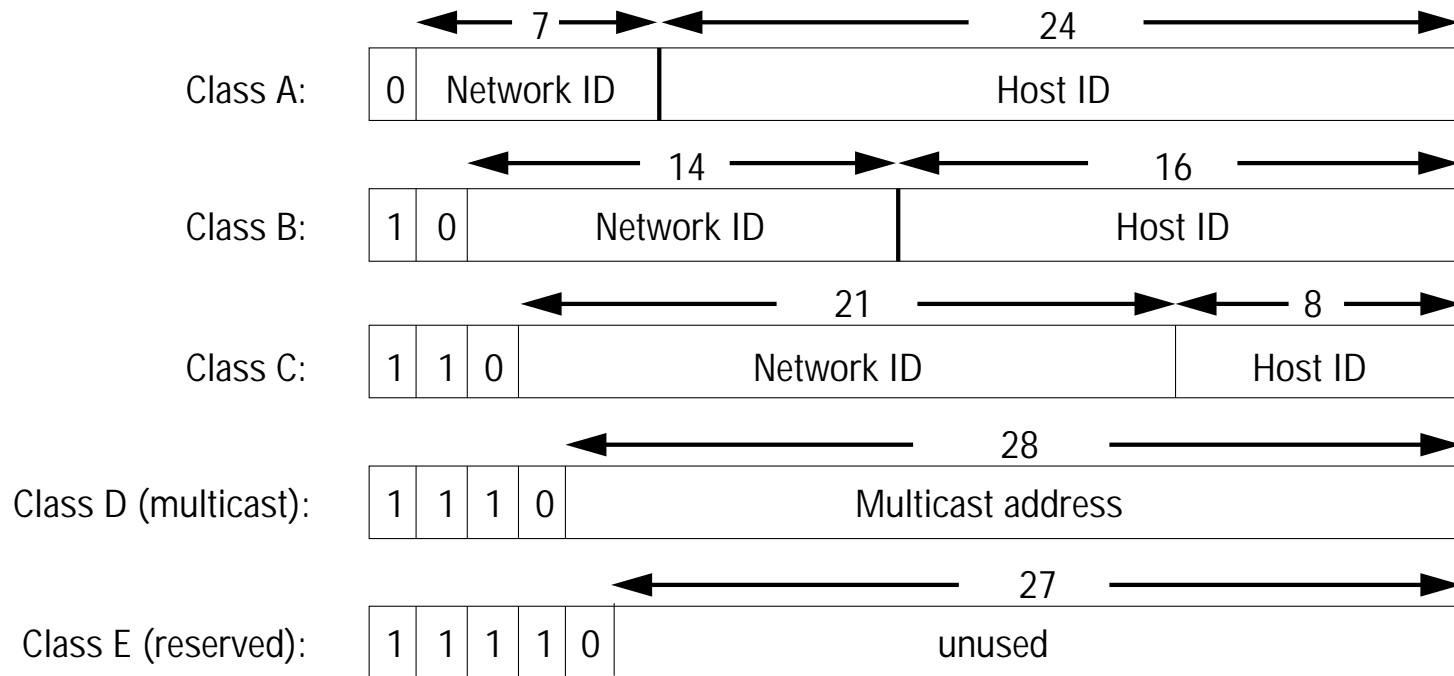


Figure 3.16 Decimal representation of Internet addresses

	<i>octet 1</i>	<i>octet 2</i>	<i>octet 3</i>		<i>Range of addresses</i>			
Class A:	<i>Network ID</i> 1 to 127	0 to 255	<i>Host ID</i> 0 to 255	0 to 255	1.0.0.0 to 127.255.255.255			
Class B:	<i>Network ID</i> 128 to 191		0 to 255	<i>Host ID</i> 0 to 255				
Class C:	192 to 223	<i>Network ID</i> 0 to 255		0 to 255	<i>Host ID</i> 1 to 254			
Class D (multicast):	<i>Multicast address</i> 224 to 239				0 to 255	0 to 255	1 to 254	224.0.0.0 to 239.255.255.255
Class E (reserved):	240 to 255	0 to 255	0 to 255	1 to 254	240.0.0.0 to 255.255.255.255			

Figure 3.17 IP packet layout

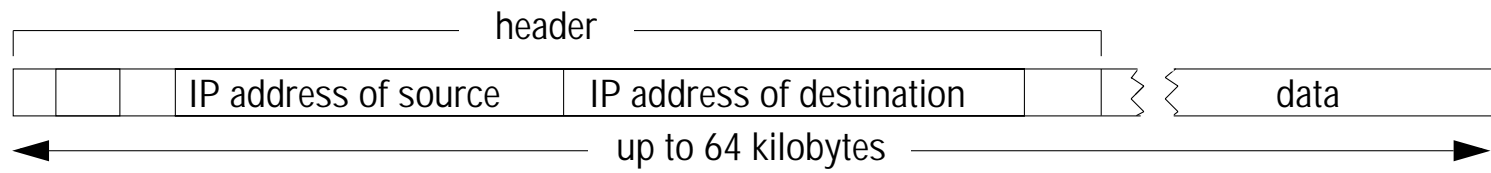


Figure 3.18 IPv6 header layout

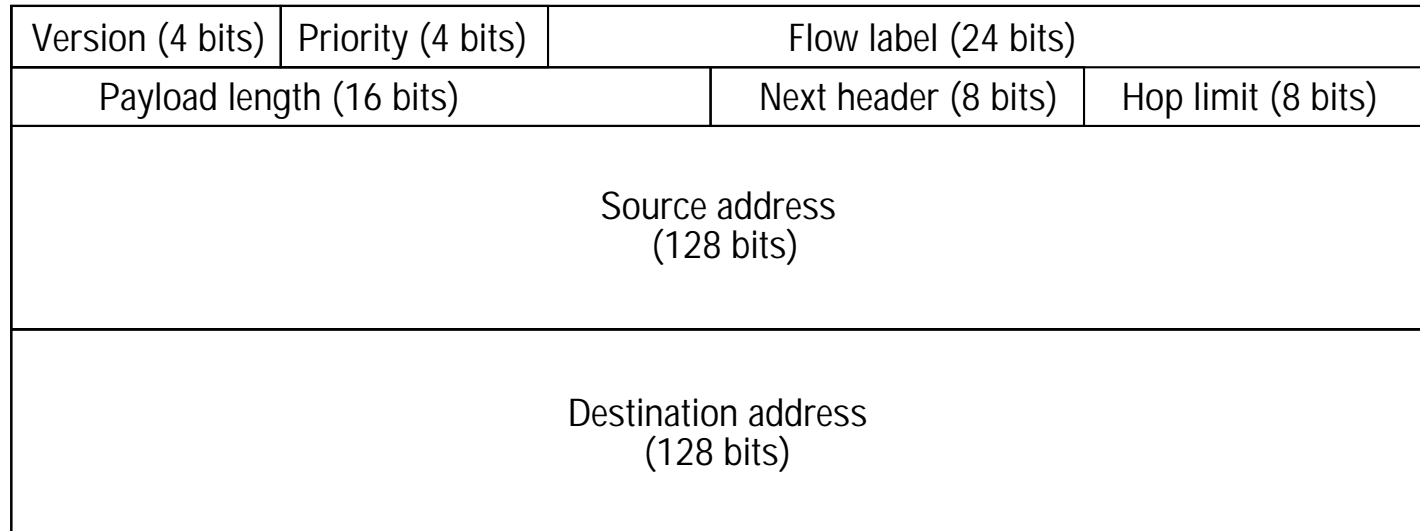


Figure 3.19 The MobileIP routing mechanism

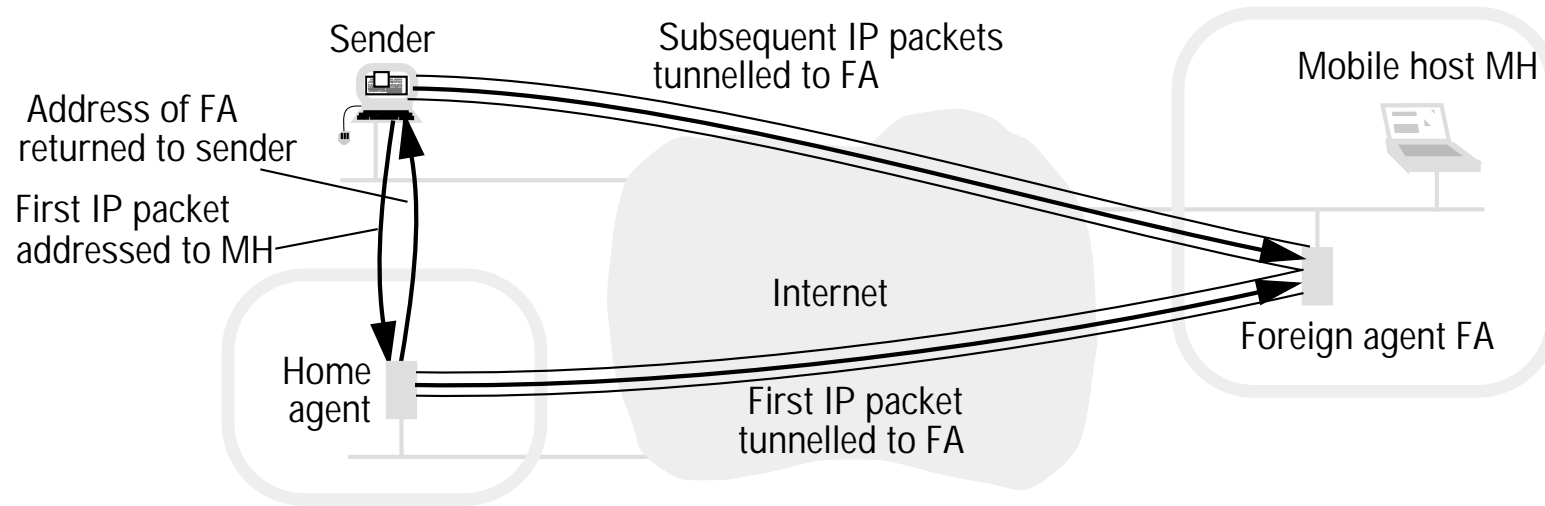
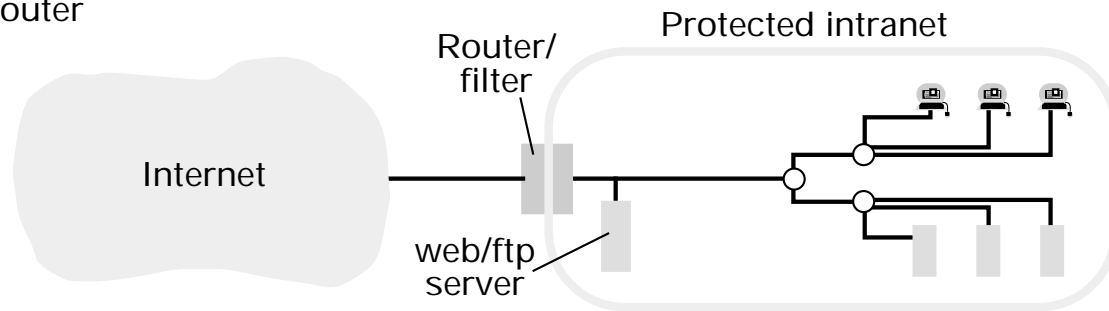
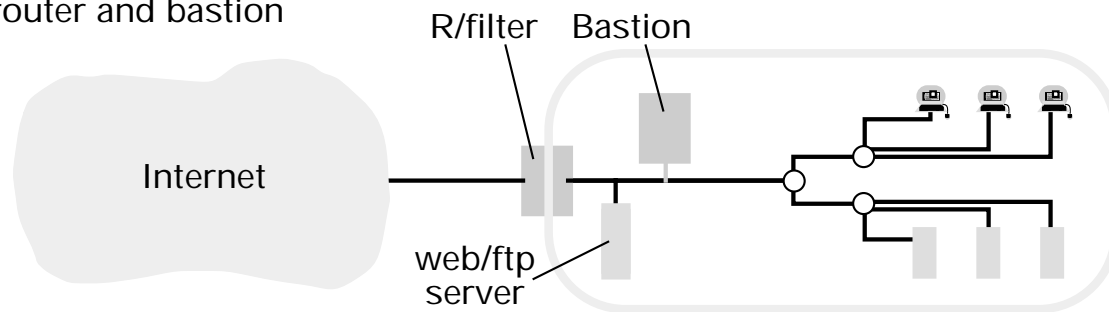


Figure 3.20 Firewall configurations

a) Filtering router



b) Filtering router and bastion



c) Screened subnet for bastion

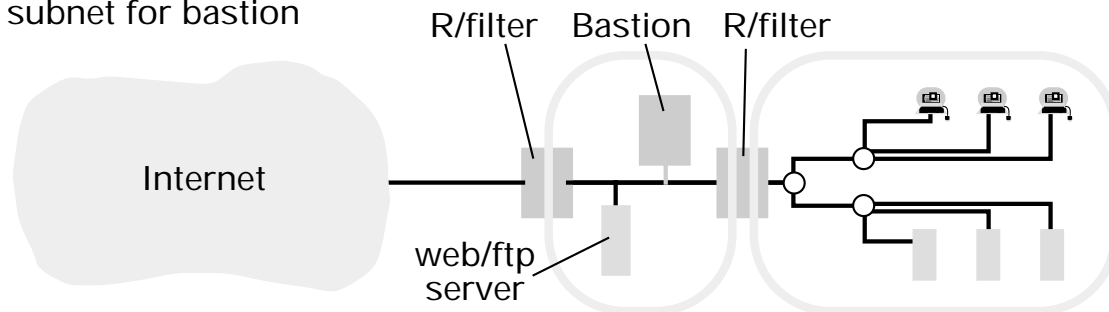


Figure 3.21 IEEE 802 network standards

<i>IEEE No.</i>	<i>Title</i>	<i>Reference</i>
802.3	CSMA/CD Networks (Ethernet)	[IEEE 1985a]
802.4	Token Bus Networks	[IEEE 1985b]
802.5	Token Ring Networks	[IEEE 1985c]
802.6	Metropolitan Area Networks	[IEEE 1994]
802.11	Wireless Local Area Networks	[IEEE 1999]

Figure 3.22 Wireless LAN configuration

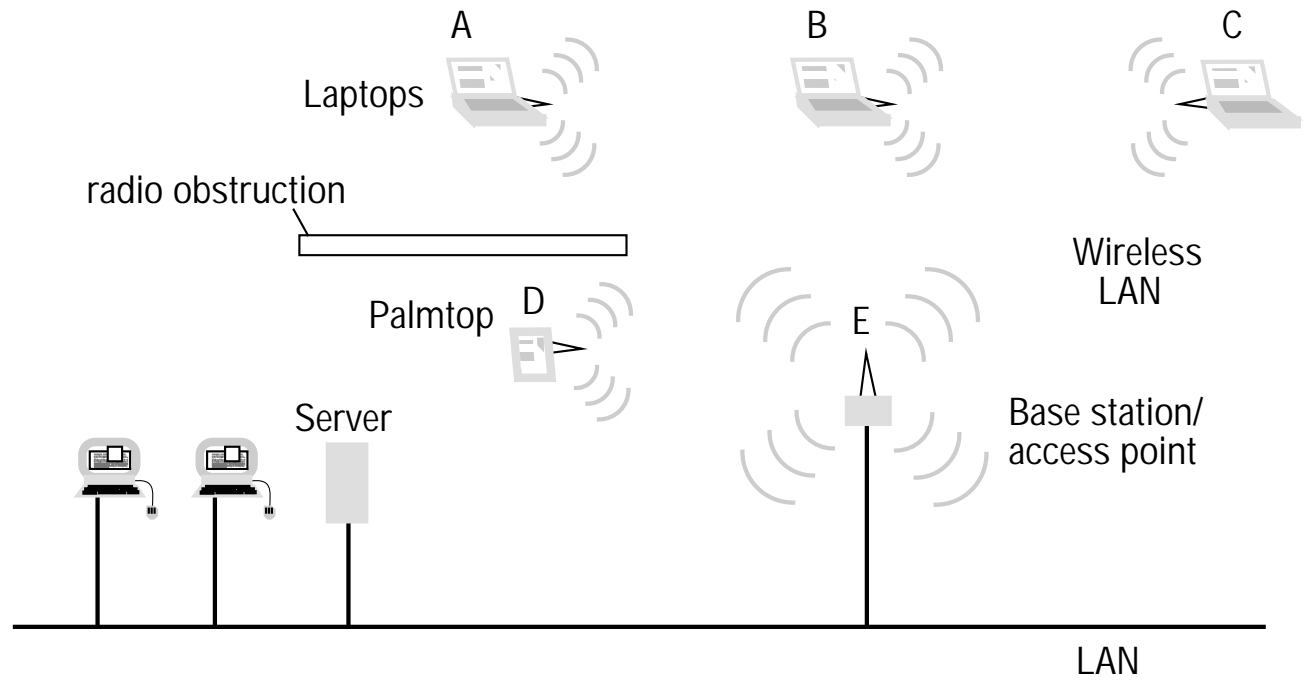


Figure 3.23 ATM protocol layers

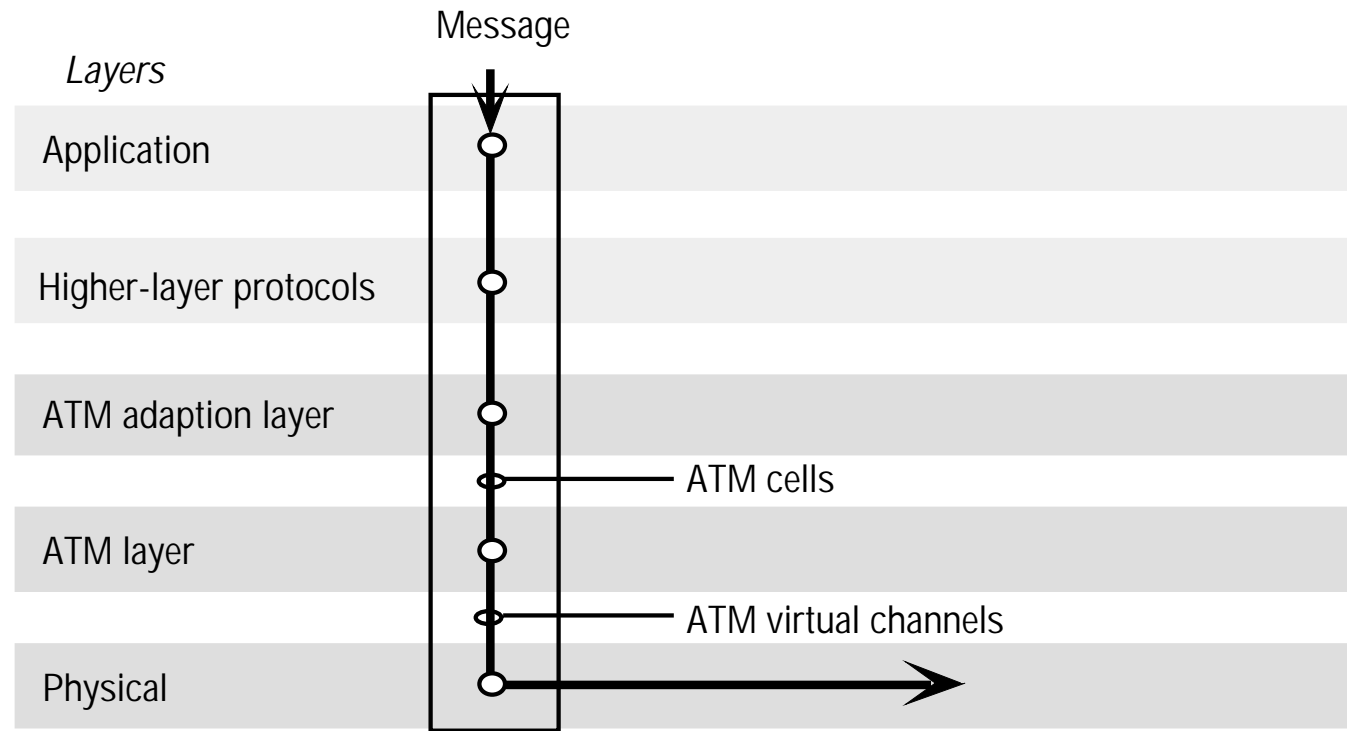


Figure 3.24 ATM cell layout

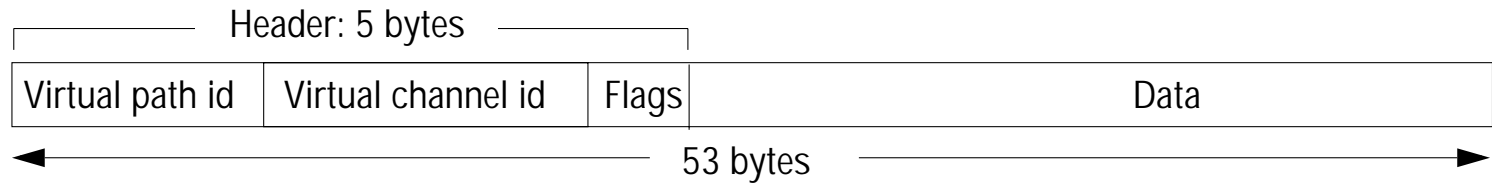


Figure 3.25 Switching virtual paths in an ATM network

